

# Some new weighted Leavitt path algebra homomorphisms

Final presentation for mathematics honours

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## Some history

- ▶ The subject of Leavitt path algebras, or Lpas, is an active research area
- ▶ It is a new subject whose seminal papers came from two groups of authors:
  - ▶ Gene Abrams and Gonzalo Aranda Pino in *The Leavitt path algebra of a graph*
  - ▶ P. Ara, M. A. Moreno, and E. Pardo in *Nonstable K-theory for Graph Algebras*
- ▶ The work in those seminal papers was initiated around the same time in 2004
- ▶ However, these groups worked independently of one another
- ▶ The genesis of the subject is in earlier work during the 1950's by W. G. Leavitt in *The Module Type of a Ring*.

## What I found difficult

- ▶ I did not have a strong background in algebra or analysis
- ▶ Lpas are defined using generators and relations
- ▶ It is not immediately clear that the presentation gives something non-zero
- ▶ In general it is difficult to find representations of Lpas in finite matrices
- ▶ It is not obvious, for example, how to know when any two given elements of an Lpa are equal
- ▶ There is, however, a representation for any Lpa inside infinite dimensional matrices over a field

## What went well?

- ▶ I wrote some software in the Haskell language to help me try to understand Lpas
- ▶ Actually I worked with weighted Lpas, or wLpas, which are a generalisation of Lpas due to Rozbeh Hazrat.
- ▶ My software can do the following things
  - ▶ Can check if two elements of an wLpa are equal
  - ▶ Caveat: I did not yet prove termination of my algorithm for checking equality!
  - ▶ Can enumerate the basis elements of a wLpa
  - ▶ Can check if a homomorphism defined on the generators of a wLpa is well-defined
  - ▶ Can compute some wLpa invariants
- ▶ I used this software to help me obtain some new homomorphisms

# The mathematics

- ▶ What are Leavitt path algebras?
- ▶ I will give a formal definition shortly
- ▶ The subject combines graph theory and ring theory
- ▶ The finite dimensional case has a more combinatorial flavour
- ▶ The general case seems to be more about ring theory

## The basics

- ▶ Let  $E$  be a directed graph
- ▶ Denote the vertices of  $E$  by  $E^0$
- ▶ Denote the edges of  $E$  by  $E^1$
- ▶ For any edge  $e$  define the *source*  $s(e)$  as the vertex at the start of the edge
- ▶ For any edge  $e$  define the *range*  $r(e)$  as the vertex at the end of the edge
- ▶ For any edge  $e$  define the “ghost-edge”  $e^*$

## The basics continued

- ▶ A *sink* is a vertex with no outgoing edges
- ▶ A *regular* vertex is a vertex that is not a sink
- ▶ Denote the set of regular vertices in a directed graph  $E$  by  $\text{Reg}(E)$
- ▶ Consider the oriented  $n$ -line graph



- ▶ Then  $v_n$  is a sink and  $v_1, \dots, v_{n-1}$  are regular.

# The formal definition

- ▶ The Leavitt path algebra on a graph is the free  $K$ -algebra generated by the vertices, edges, and ghost-edges, modulo the following relations

**V**  $uv = \delta_{uv}v$  for all  $u, v \in E^0$

**E1**  $s(e)e = er(e) = e$  for all  $e \in E^1$

**E2**  $r(e)e^* = e^*s(e) = e^*$  for all  $e^* \in (E^1)^*$

**CK1**  $e^*f = \delta_{ef}r(e)$  for all  $e, f \in E^1$

**CK2**  $v = \sum_{\substack{e \in E^1 \\ s(e)=v}} ee^*$  for all  $v \in \text{Reg}(E)$

- ▶ What is a free  $K$ -algebra?

- ▶ It is like a non-commutative version of a polynomial ring with coefficients from a field  $K$

## The formal definition continued

- ▶ What does modulo mean?
- ▶ It means quotient
- ▶ What is a quotient?
  - ▶ A quotient is a way to obtain a new algebra from an old one by including some relations
  - ▶ It is defined in most algebra texts using set theory.

## The motivation

- ▶ Why these relations and what do they mean?
- ▶ In the case of finite directed acyclic graphs these relations describe an algebra with some interesting combinatorial properties
- ▶ What does acyclic mean?
- ▶ It means there are no cycles
- ▶ What is a cycle?
- ▶ A *cycle* in a directed graph is a path  $e_1, \dots, e_n$  such that  $s(e_i) \neq s(e_j)$  for every  $i \neq j$ , and  $r(e_n) = s(e_1)$

## The motivation continued

- ▶ What is a path?
- ▶ A *path* in a directed graph is a finite sequence of edges  $e_1, \dots, e_n$ , such that  $r(e_i) = s(e_{i+1})$  for all  $1 \leq i \leq n - 1$ .
- ▶ Note: If we are discussing the Lpa of a graph, then the words path and cycle refer to the corresponding products of these edges, rather than merely sequences of edges

## The oriented $n$ -line graph

- Let us revisit the oriented  $n$ -line graph example,



- We want to be able to add, multiply, and subtract vertices, edges, and ghost-edges
- How do we define this in a meaningful way?
- Use matrices
- We represent the vertex  $v_1$  as the matrix

$$\begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ 0 & 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 \end{bmatrix}$$

## The oriented $n$ -line graph continued

- ▶ We represent the vertex  $v_2$  as the matrix

$$\begin{bmatrix} 0 & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 \end{bmatrix}$$

- ▶ See the pattern?
- ▶ Represent vertex  $v_k$  as the single-entry matrix with 1 in the  $k$ th diagonal entry
- ▶ What about edges?

## The oriented $n$ -line graph continued

Representations for the edges are as follows,

- We represent the edge  $e_1$  as the matrix

$$\begin{bmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 \end{bmatrix}$$

- We represent the edge  $e_2$  as the matrix

$$\begin{bmatrix} 0 & 0 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 \end{bmatrix}$$

## The oriented $n$ -line graph continued

- ▶ We represent the edge  $e_k$  as the single-entry matrix with 1 in the  $k$ th row and  $(k + 1)$ th column
- ▶ What about ghost-edges?
- ▶ The  $*$  operation corresponds to the matrix transpose or adjoint
- ▶ For example  $e_1^*$  is represented by

$$\begin{bmatrix} 0 & 0 & 0 & \dots & 0 \\ 1 & 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 \end{bmatrix}$$

- ▶ The ghost-edge  $e_k^*$  is represented by the single-entry matrix with 1 in the  $(k + 1)$ th row and the  $k$ th column

## The oriented $n$ -line graph continued

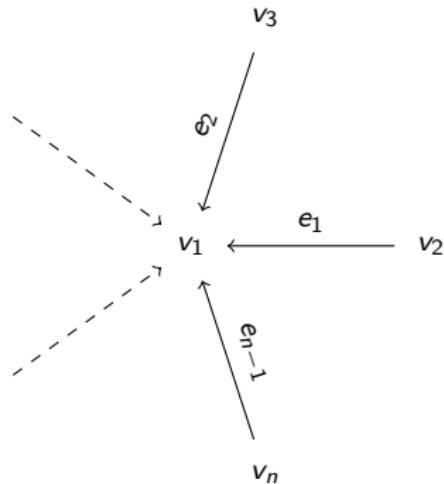
- ▶ Let  $M_K(n)$  be the  $K$ -algebra of square  $n$  by  $n$  dimensional matrices over the field  $K$
- ▶ For any directed graph  $E$  we will denote the Leavitt path algebra of  $E$  with coefficients from the field  $K$  by  $L_K(E)$
- ▶ For the oriented  $n$ -line graph,  $A_n$ , we get  $L_K(A_n) \cong M_K(n)$
- ▶ What about other directed acyclic graphs?

## Directed acyclic graphs

- ▶ Consider any directed acyclic graph  $E$  with a sink  $v$
- ▶ Since the graph is finite acyclic, there must be a finite number of paths that end at  $v$ . Let this number be  $n(v)$ .
- ▶ The ideal generated by  $v$ ,  $I(v)$  is isomorphic to  $M_K(n(v))$
- ▶ If there is only one sink, then this ideal is the entire Lpa
- ▶ If there is more than one sink, then the Lpa is given by the direct product of the ideals generated by the sinks.

## Directed acyclic graphs continued

- ▶ Consider this example,



- ▶ Denote the above graph with  $n$  vertices as  $B_n$
- ▶ Denote the oriented  $n$ -line graph by  $A_n$
- ▶ Then  $L_K(B_n) \cong L_K(A_n) \cong M_K(n)$

## Directed acyclic graphs continued

- ▶ For any  $n > 0$  there are a finite number of directed acyclic graphs whose Lpa is isomorphic to  $M_K(n)$
- ▶ Counting these graphs appears to be a non-trivial problem
- ▶ This problem seems to be worthy of further study
- ▶ I was not able to find anything in the literature about this so far

## Graphs with cycles

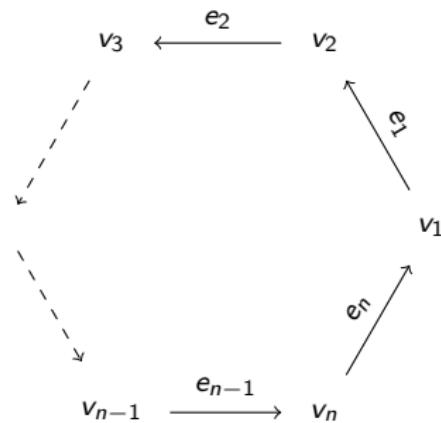
- ▶ For certain types of graphs with cycles, we may still be able to obtain representations in finite matrices over Laurent polynomial rings
- ▶ Consider the loop graph



- ▶ The Lpa for this graph is isomorphic to  $K[x, x^{-1}]$ , the Laurent polynomial ring over  $K$

## Graphs with cycles continued

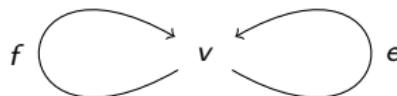
- Or the  $n$ -cycle graph



- The Lpa for this graph is isomorphic to  $M_{K[x,x^{-1}]}(n)$

## Graphs with cycles continued

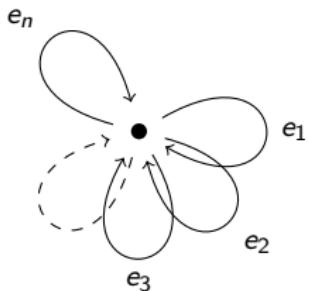
- ▶ Consider the 2-petal graph



- ▶ The Lpa of this graph, call it  $R$ , has an unusual property
- ▶ As left  $R$ -modules, we have  $R \cong R \oplus R$
- ▶ A *module* is like a vector space over a ring
- ▶ W. G. Leavitt was the first person to produce examples of rings having the above property

## Graphs with cycles continued

- More generally, consider the  $n$ -petal graph for some fixed integer  $n > 0$ ,



- Denote by  $R^n$  the direct sum  $\underbrace{R \oplus \cdots \oplus R}_{n \text{ times}}$
- The Lpa of this graph, call it  $R$ , satisfies  $R \cong R^n$
- Moreover this  $n$  is the smallest integer  $m$  such that  $R \cong R^m$

# Equality

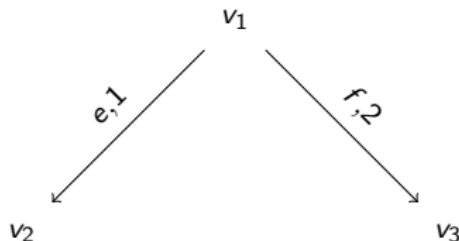
- ▶ In the general case of an Lpa, how do we know if two given elements of the algebra are equal?
- ▶ This is a non-trivial issue since we may not necessarily have a finite representation for the algebra
- ▶ There is a theorem that gives a *basis* for any Lpa, viewed as a  $K$ -module
- ▶ Checking equality amounts to reducing elements to their basis form and comparing them
- ▶ This basis will only be finite if the graph is finite acyclic

## A Haskell implementation

- ▶ I coded an algorithm using the Haskell language to check for equality of  $wLpa$  elements
- ▶ The algorithm works as follows,
  1. First convert the given element of the  $wLpa$  to normal form. That is, a  $K$ -linear combination of paths
  2. For each path in this expression, find any *forbidden* sub-paths. The term *forbidden* is a technical term!
  3. For each forbidden sub-path found, rewrite this sub-path using the  $wLpa$  relations to remove it
  4. Go back to the start and iterate this procedure until no forbidden sub-paths remain. Then stop.
- ▶ A path that contains no forbidden sub-paths is called a *nod-path*
- ▶ Proving termination of the above algorithm would show that nod-paths are a spanning set for the  $wLpa$

## Weighted graphs

- ▶ We assign positive integer weights to the edges of a directed graph. For example  $(E, w)$  given as,



- ▶ Here,  $E$  is the unweighted directed graph and  $w$  is a function that maps  $E^{st} \rightarrow \mathbb{Z}^+$
- ▶ So in this example  $w(e) = 1$  and  $w(f) = 2$
- ▶ The vertex set is given by

$$E^0 = \{v_1, v_2, v_3\}$$

- ▶ The edge set is given by

$$E^1 = \{e_1, f_1, f_2\}$$

## Weighted graphs continued

The formal definition is as follows,

- ▶ Let  $E^0$  be a countable set called *vertices*
- ▶ Let  $E^{st}$  be a countable set called *structured edges*
- ▶ Let there be maps  $s, r : E^{st} \rightarrow E^0$
- ▶ Let there be a *weight map*  $w : E^{st} \rightarrow \mathbb{Z}^+$
- ▶ Define the *edges* as the set  $E^1 = \bigcup_{\alpha \in E^{st}} \{\alpha_i \mid 1 \leq i \leq w(\alpha)\}$

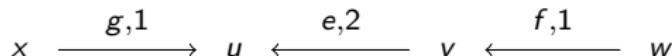
# Weighted Leavitt path algebras

The formal definition,

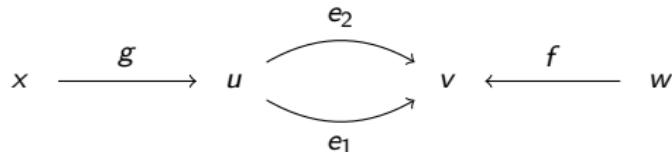
- ▶ A wLpa is the free  $K$ -algebra on  $E^0 \cup E^1 \cup (E^1)^*$ , subject to the following relations
  1.  $v_i v_j = \delta_{ij} v_i$  for every  $v_i, v_j \in E^0$ .
  2.  $s(\alpha)\alpha_i = \alpha_i r(\alpha) = \alpha_i$  and  $r(\alpha)\alpha_i^* = \alpha_i^* s(\alpha) = \alpha_i^*$  for all  $\alpha \in E^{st}$  and  $1 \leq i \leq w(\alpha)$ .
  3.  $\sum_{\{\alpha \in E^{st} | s(\alpha) = v\}} \alpha_i \alpha_j^* = \delta_{ij} v$  for fixed  $1 \leq i, j \leq \max\{w(\alpha) | \alpha \in E^{st}, s(\alpha) = v\}$ , for all  $v \in E^0$ .
  4.  $\sum_{1 \leq i \leq \max\{w(\alpha), w(\alpha')\}} \alpha_i^* \alpha'_i = \delta_{\alpha\alpha'} r(\alpha)$ , for all  $\alpha, \alpha' \in E^{st}$ .
- ▶ In the above, if  $e \in E^{st}$  is a structured edge and  $i > w(e)$  then  $e_i = 0$ .

## Some wLpa examples

- ▶ Many simple examples of wLpas are isomorphic to unweighted Lpas
- ▶ For example the wLpa of



- ▶ is isomorphic to the Lpa of



## Some wLpa examples continued

The isomorphism is given on the generators as,

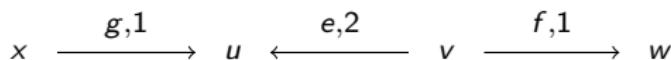
$$\begin{array}{ll} u \mapsto u & e_1 \mapsto e_1^* \\ v \mapsto v & e_2 \mapsto e_2^* \\ w \mapsto w & f_1 \mapsto f \\ x \mapsto x & g_1 \mapsto g \end{array}$$

## Some wLpa examples continued

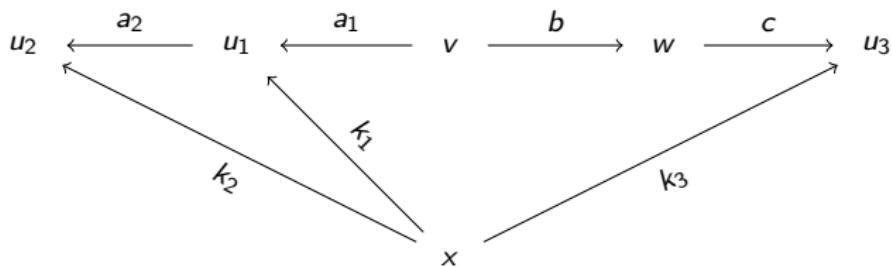
- ▶ To see this is a well defined homomorphism,
  1. Start with the free algebra on the vertices, edges, and ghost edges of the weighted graph
  2. As soon as you have a map from the generators to a ring, it induces a homomorphism on the level of algebra, by the properties of a free algebra
  3. Next one needs to check that the weighted relations will be mapped to zero in the Lpa.
  4. This guarantees a map from the quotient of the free algebra modulo the weighted relations into the Lpa
- ▶ Step 3 can be checked by computer.

## Some wLpa examples continued

- ▶ The wLpa of



- ▶ is isomorphic to the Lpa of



## Some wLpa examples continued

The isomorphism is given on the generators as,

$$u \mapsto u_1 + u_2 + u_3$$

$$e_1 \mapsto a_1$$

$$v \mapsto v$$

$$e_2 \mapsto a_1 a_2 + bc$$

$$w \mapsto w$$

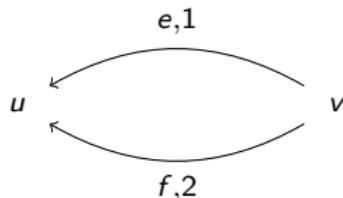
$$f_1 \mapsto b$$

$$x \mapsto x$$

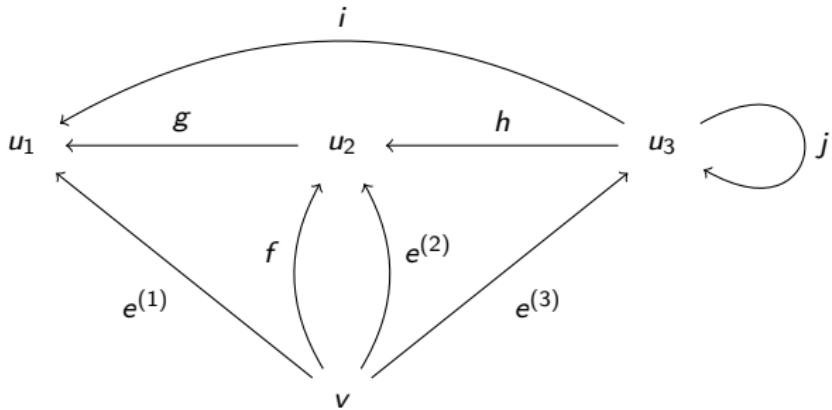
$$g_1 \mapsto k_1 + k_2 + k_3$$

## Some wLpa examples continued

- ▶ A more complicated example. The wLpa of



- ▶ this is isomorphic to the Lpa of



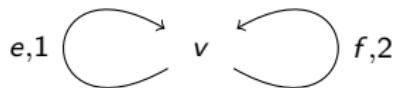
## Some wLpa examples continued

The isomorphism is given on the generators as,

$$\begin{array}{ll} u \mapsto u_1 + u_2 + u_3 & e_1 \mapsto e_1 + e_2 + e_3 \\ v \mapsto v & f_1 \mapsto f \\ & f_2 \mapsto fg + e_1 i^* + e_2 h^* + e_3 j^* \end{array}$$

## An unsolved problem

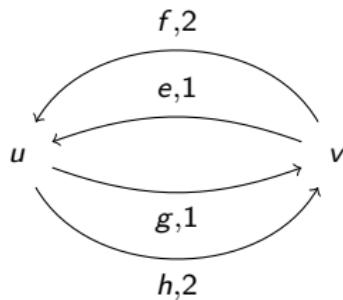
- ▶ The previous example is from a paper in arXiv entitled *The V-monoid of a weighted Leavitt path algebra* by Raimund Preusser
- ▶ In the same paper Preusser gives the following example



- ▶ At the present time it is unknown whether this wLpa is isomorphic to an unweighted Lpa

## Some new homomorphisms

- ▶ In my honours thesis I introduced the following weighted graph,



- ▶ I was able to obtain some new homomorphisms using this

## Some new homomorphisms continued

$$L_K \left( \begin{matrix} e,1 & \curvearrowright & v & \curvearrowleft & f,2 \end{matrix} \right) \rightarrow L_K \left( \begin{matrix} u & \xleftarrow{e,1} & f,2 & \xrightarrow{g,1} & v \\ & \xleftarrow{h,2} & & & \end{matrix} \right)$$

$$v \mapsto v + u$$

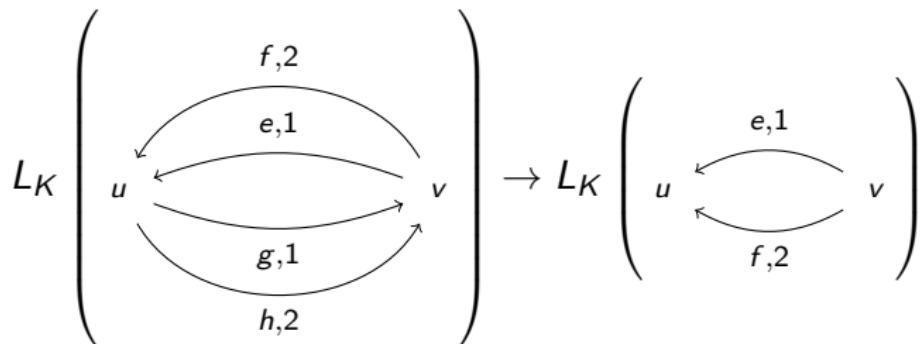
$$f_1 \mapsto f_1 + h_1$$

$$e_1 \mapsto e_1 + g_1$$

$$f_2 \mapsto f_2 + h_2$$

- ▶ Call this map  $\eta$
- ▶ I claim that  $\eta$  is a monomorphism.

## Some new homomorphisms continued



$$u \mapsto u$$

$$e_1 \mapsto e_1$$

$$f_2 \mapsto f_2$$

$$h_1 \mapsto f_1^*$$

$$v \mapsto v$$

$$f_1 \mapsto f_1$$

$$g_1 \mapsto f_2^*$$

$$h_2 \mapsto e_1^*$$

- ▶ Call this map  $\rho$
- ▶ Observe that the generators of the codomain all appear in the range of  $\rho$
- ▶ Therefore  $\rho$  is an epimorphism

## Some new homomorphisms continued

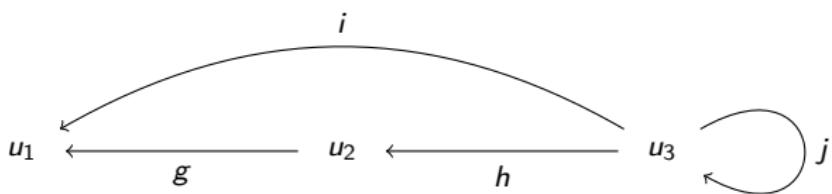
- ▶ Define  $Q = \{h_2^* - e_1, h_1^* - f_1, g_1^* - f_2\}$
- ▶ In my honours thesis I show that  $\ker(\rho) = I(Q)$
- ▶ Let  $\theta = \rho \circ \eta$
- ▶ So  $\theta$  is given by

$$\begin{array}{ll} v \mapsto v + u & f_1 \mapsto f_1 + f_1^* \\ e_1 \mapsto e_1 + f_2^* & f_2 \mapsto f_2 + e_1^* \end{array}$$

- ▶ Let  $R = \{e_1 - f_2^*, f_1 - f_1^*\}$
- ▶ I claim that  $\ker(\theta) = I(R)$

## Some new homomorphisms continued

- ▶ In my honours thesis I introduced the following unweighted graph



- ▶ I was able to obtain some new homomorphisms using this

## Some new homomorphisms continued

There is a homomorphism,

$$L_K \left( \begin{array}{ccc} & i & \\ u_1 & \xleftarrow{g} & u_2 & \xleftarrow{h} & u_3 \\ & j & \end{array} \right) \longrightarrow$$

$$L_K \left( \begin{array}{ccc} e,1 & \xrightarrow{\quad} & v & \xleftarrow{\quad} & f,2 \end{array} \right)$$

## Some new homomorphisms continued

- ▶ It is given on the generators by

$$u_1 \mapsto x^*x$$

$$u_2 \mapsto xx^*$$

$$u_3 \mapsto y^*y - x^*x - xx^*$$

$$g \mapsto x$$

$$h \mapsto y - yz$$

$$i \mapsto y(z - yy^*)$$

$$j \mapsto y^2y^*$$

where,

$$x = f_1^*f_2 \quad y = f_2^*e_1 \quad z = f_2^*f_2$$

- ▶ I claim this map is mono

## Some new homomorphisms continued

There is a representation in 2 by 2 matrices as follows,

$$L_K \left( \begin{array}{cc} & e,1 \\ u & \swarrow \searrow \\ & f,2 \\ & v \end{array} \right) \hookrightarrow M_R(2)$$

where

$$R = L_K \left( \begin{array}{ccccc} & & & & \\ & e,1 & \curvearrowleft & v & \curvearrowleft \\ & & & & f,2 \end{array} \right)$$

## Some new homomorphisms continued

It is given on the generators as,

$$e_1 \mapsto \begin{bmatrix} 0 & e_1 \\ 0 & 0 \end{bmatrix}$$

$$f_1 \mapsto \begin{bmatrix} 0 & f_1 \\ 0 & 0 \end{bmatrix}$$

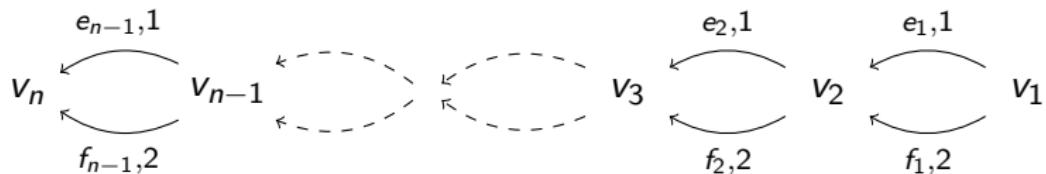
$$f_2 \mapsto \begin{bmatrix} 0 & f_2 \\ 0 & 0 \end{bmatrix}$$

$$u \mapsto \begin{bmatrix} 0 & 0 \\ 0 & v \end{bmatrix}$$

$$v \mapsto \begin{bmatrix} v & 0 \\ 0 & 0 \end{bmatrix}$$

## Some new homomorphisms continued

- More generally, let  $(G_n, w_{G_n})$  be the following weighted graph,



- There is a representation,

$$L_K(G_n, w_{G_n}) \hookrightarrow M_R(n)$$

for all  $n > 1$

## Some new homomorphisms continued

- ▶ It is given by,

$$\begin{aligned} e_{i,1} &\mapsto e_1 c_{i,i+1} & f_{i,1} &\mapsto f_1 c_{i,i+1} & f_{i,2} &\mapsto f_2 c_{i,i+1} \\ v_i &\mapsto v c_{i,i} \end{aligned}$$

where  $c_{i,j}$  is the  $n$  by  $n$  single-entry matrix with 1 in the  $i$ th row and  $j$ th column.

- ▶ Question: how to generalise this result further?

## Some new homomorphisms continued

- ▶ There is a homomorphism as follows,

$$L_K \left( \begin{array}{ccc} & f,2 & \\ & \swarrow \curvearrowright & \searrow \\ u & e,1 & v \\ & \swarrow \curvearrowright & \searrow \\ & g,1 & \\ & \swarrow \curvearrowright & \searrow \\ & h,2 & \end{array} \right) \rightarrow L_{K[x,x^{-1}]} \left( \begin{array}{ccc} & e,1 & \\ & \swarrow \curvearrowright & \searrow \\ u & & v \\ & \swarrow \curvearrowright & \searrow \\ & f,2 & \end{array} \right)$$

- ▶ Note that in the codomain the coefficients come from a Laurent polynomial  $*$ -ring
- ▶ We assume that  $x^* = x^{-1}$

## Some new homomorphisms continued

- ▶ It is given on the generators as,

$$\begin{array}{ll} e_1 \mapsto xe_1 & g_1 \mapsto xf_2^* \\ f_1 \mapsto xf_1 & h_1 \mapsto xf_1^* \\ f_2 \mapsto xf_2 & h_2 \mapsto xe_1^* \end{array} \quad \begin{array}{l} u \mapsto u \\ v \mapsto v \end{array}$$

- ▶ I claim this map is mono

## Some new homomorphisms continued

- ▶ Composing an earlier example with the above example gives a monomorphism,

$$L_K \left( \begin{array}{ccc} & e,1 & \\ & \curvearrowleft & \curvearrowright \\ & v & f,2 \end{array} \right) \rightarrow L_{K[x,x^{-1}]} \left( \begin{array}{ccc} & e,1 & \\ & \curvearrowleft & \curvearrowright \\ u & & v \\ & \curvearrowleft & \curvearrowright \\ & f,2 & \end{array} \right)$$

- ▶ The map is given by,

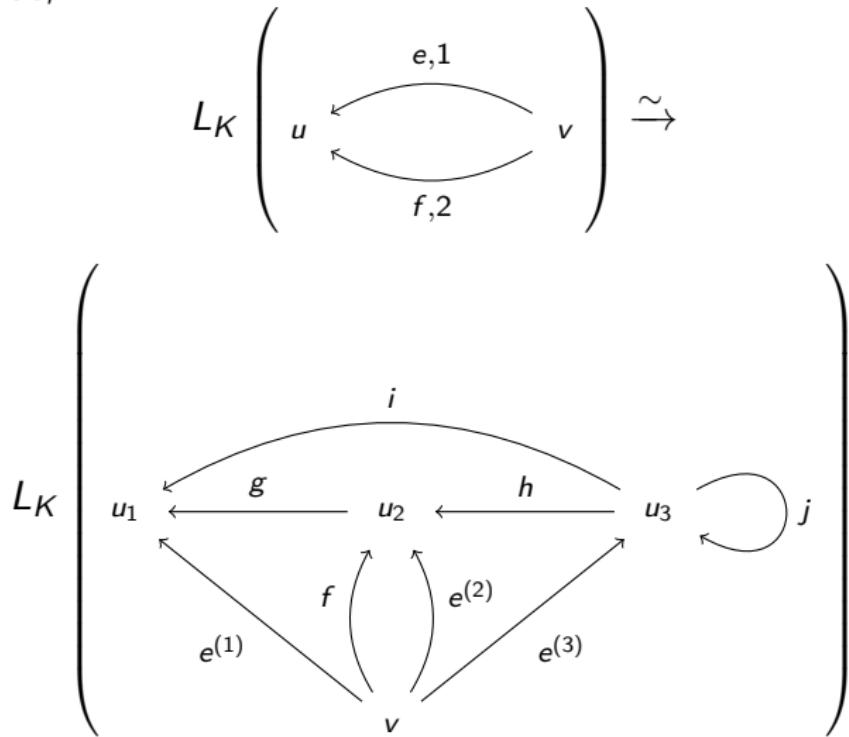
$$e_1 \mapsto x(e_1 + f_2^*)$$

$$f_1 \mapsto x(f_1 + f_1^*) \quad v \mapsto u + v$$

$$f_2 \mapsto x(f_2 + e_1^*)$$

## Some new homomorphisms continued

Recall that,



## Some new homomorphisms continued

- ▶ It follows that our  $wLpa$ ,

$$L_K \left( \begin{array}{ccc} e,1 & \curvearrowright & f,2 \\ & v & \curvearrowleft \end{array} \right)$$

- ▶ is sitting inside an unweighted  $Lpa$  over a Laurent polynomial  $*$ -ring

## Where is the code?

- ▶ The Haskell code I used is freely available on Github at  
[https://github.com/rzil/honours/tree/master/  
LeavittPathAlgebras](https://github.com/rzil/honours/tree/master/LeavittPathAlgebras)
- ▶ All of the above homomorphisms can be verified computationally using this code
- ▶ Coming up with these maps, I found enjoyable. I would encourage you to try it!

## Further work

- ▶ In the above mentioned examples, there were a number of claims yet to receive proof
- ▶ For example, showing that certain maps are mono or finding generating sets for the kernels
- ▶ One idea for showing some of the injectivity claims, suggested by R. Preusser, is to show that nod-paths are mapped to distinct nod-paths
- ▶ Most of the maps given above were verified using computer. I would like to prove some of them on paper.
- ▶ The basis algorithm described earlier has not been proven to terminate. I would like to prove this.

Thank you for listening!